

Learning repetition and optionality

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Human language acquisition

Learning

Blah blah...



Blah
blah...

Photograph by Andrew Hetherington, Scientific American
July 20 2011

<http://www.scientificamerican.com/article.cfm?id=heparich-baby-brains-signal-later-language-problems>

Human language acquisition

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Blah blah...



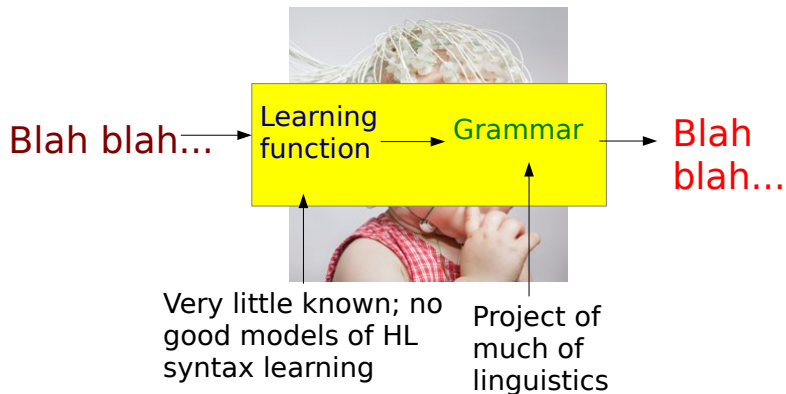
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Human language acquisition

Learning



Overview

The phenomenon In human language, adjuncts are **optional** and often **repeatable**

The question How do formal learning models handle them?

Adjuncts

Generally adjectives, adverbs, prepositional phrases

- (1) a. My love is like a rose.
b. My love is like a red red rose.
- (2) a. I'm tired!
b. I'm really really really really tired!
- (3) He suddenly (*suddenly suddenly) smiled.

Learners we'll look at today

- 1 0-reversible learner (Angluin, 1982)
- 2 Substitutable CFGs (Clark, 2010)
- 3 PAC learner for PDFAs (Clark and Thollard, 2004)
- 4 (n-gram learner (García and Vidal, 1990))

Findings

- ① 0-reversible learner (Angluin, 1982)
 - Optionality \leftrightarrow Repetition
 - one repetition \rightarrow indefinite repetition
 - need $ux^n v, ux^{n+1} v$
- ② Substitutable CFGs (Clark, 2010)
 - repetition \leftrightarrow optionality
 - one repetition \rightarrow indefinite repetition
 - need $ux^n v, ux^{n+1} v$
- ③ PAC learner for PDFAs (Clark and Thollard, 2004)
 - Representative sample \rightarrow learns repetition
 - Optionality $\not\leftrightarrow$ Repetition
- ④ (n-gram learner (García and Vidal, 1990))
 - Given up to n x 's in a row in context (u, v) , generalises to $ux^* v \subseteq L$
 - No generalisation from optionality to repetition or vice-versa

Learnability

A very weak claim *For some definition of “learn” and some definition of “language”, humans learn language*

Learnability

Definition (Language)

Given a finite set Σ , Σ^* is the set of all finite sequences of elements of Σ .
 L is a language iff $L \subseteq \Sigma^*$

Definition (Learner)

A function from texts (samples from L) to grammars

Definition (Learn)

A learner learns a class of languages if it distinguishes them from each other

Learnability – a note

- A learner makes assumptions about the input strings and encodes them into its hypothesis grammar.
- So even if the input to a class X learner is not generated by a language of class X , the grammar that the learner hypothesises will be of class X .
- Thus the learner only correctly learns the language if it was from that class in the first place.
- This is what it means for a learner to learn a class of languages

Learnability

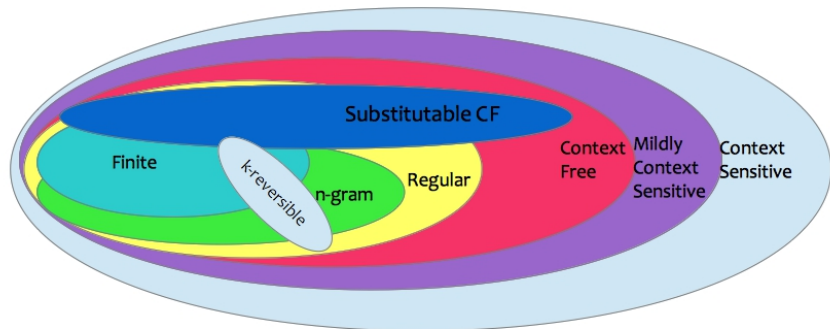
Definition (Gold learning)

(a.k.a. Identification in the limit from positive data) A language class is Gold-learnable if there is a function that will ultimately correctly converge on a grammar for every language in the class (Gold, 1967)

Definition (PAC learning)

$\forall 0 < \epsilon < 0.5, 0 < \delta < 0.5$ a Probably Approximately Correct learner outputs hypothesis grammars that are, with probability $1 - \delta$, ϵ -close to correct. (Valiant, 1984)

Chomsky hierarchy



Optionality and repetition

Definition (Optional)

$x \in \Sigma^*$ is *optional in context* u,v iff $uv \in L$ and $uxv \in L$

Definition (Repeatable)

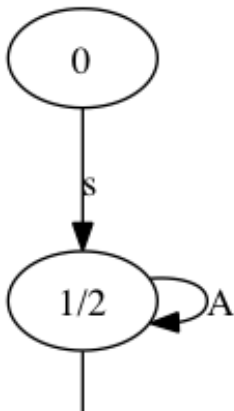
$x \in \Sigma^*$ is *repeatable in context* u,v iff $ux^+v \subseteq L$

0-reversible learner

Definition

A FSA is 0-reversible iff it is deterministic both forward and backward

If L is 0-reversible then for all strings u, v , if u and v share one suffix, they share all suffixes.

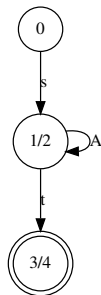
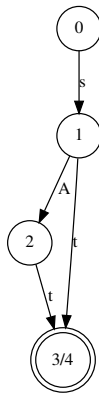
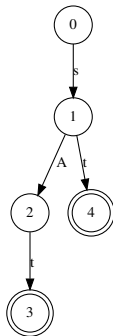


$L = sA^*t$
 $s, sA, sAA, sAAA\dots$
 all have suffixes
 $t, At, AAt, AAAAt\dots$

0-reversible learner

Optionality \rightarrow RepetitionSample: st, sAt

Prefix	Suffix
ϵ	st, sAt
s	t, At
st	ϵ
sA	t
sAt	ϵ



0-reversible learner

Theorem (Optionality \rightarrow Repetition)

Let $u, v, x \in \Sigma^*$ and let uv, uxv be in the sample of L . Then $ux^*v \subseteq L$

Proof.

ux and u share suffix v .

u also has suffix xv

$\rightarrow ux$ also has suffix xv ,

$\rightarrow uxxv \in L$.

uxx has suffix v as well

$\rightarrow uxx$ also must have suffix xv ,

$\rightarrow uxxxv \in L$

etc.

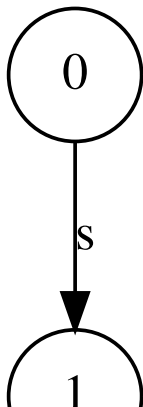


uv uxv

uxv $\rightarrow uxxv$

0-reversible learner – repetition

Sample: sAt , $sAtt$



0-reversible learner: repetition \rightarrow optionalityTheorem (repetition \rightarrow optionality)

Let $ux^k v, ux^{k+1} \in L$ for some $k > 0$. Then $uv \in L$.

Proof.

$ux^k v, ux^{k+1} \in L$

u^{k-1}, ux^k share suffix xv

ux^k also has suffix v

\rightarrow so does ux^{k-1}

$\rightarrow ux^{k-1}v \in L$

etc.



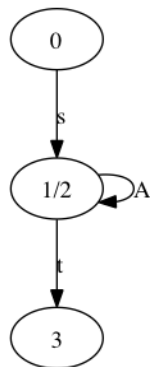
0-reversible learner - summary

- Optionality \leftrightarrow Repetition
- one repetition \rightarrow indefinite repetition

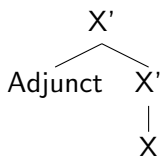
Interlude: States and repetition

- 0-reversible demonstrates something we want to capture: the notion that we should expect optionality and repetition of x in context C just in case *it doesn't matter whether x occurs in C or not*
- i.e. we're in the same state regardless of x 's presence
- → **We expect repetition and optionality to pattern together**

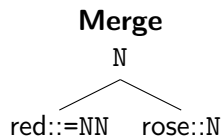
States and repetition



FSA



X-bar tree



Minimalist Grammar
 Stabler (1997)
 Kobele et al. (2007)

Substitutable Context Free learner

- CF equivalent of 0-reversible
- Learnable!

Definition (Substitutable context free language)

L is SCF iff for all $u, v, s, t, x_1, x_2 \in \Sigma^*$, if

$ux_1v \in L$ and

$ux_2v \in L$ and

$sx_1t \in L$ then

$sx_2t \in L$

i.e if two strings share one **context**, they share all contexts

Substitutable Context Free learner

Theorem (Optionality \rightarrow Repetition)

Let $u, v, x \in \Sigma^*$ and $uv, uxv \in L$ Then $ux^*v \subseteq L(G_i)$.

Proof.

By induction on the number of x s.

u, ux share context (ϵ, v) .

u also has context (ϵ, xv) so ux also must have this context.

Therefore $uxxv \in L(G_i)$.

Suppose $ux^k v \in L$. Then ux^{k-1}, ux^k share context (ϵ, v) .

ux^{k-1} also has context (ϵ, xv) so $u^k x$ also must have this context.

Therefore $ux^{k+1}v \in L(G_i)$.



$uv \quad uxv$
 $uxv \quad \rightarrow \quad uxxv$

Substitutable Context Free learner

Theorem (Repetition \rightarrow Optionality)

Let $u, v, x \in \Sigma^*$ and $ux^n v, ux^{n+1} v \in T[i]$ Then $uv \subseteq L(G_i)$.

Proof.

By induction on the number of x s. ux^{n-1}, ux^n share context (ϵ, xv) . ux^n also has context (ϵ, v) so ux^{n-1} also must have this context. Therefore $ux^{n-1}v \in L(G_i)$. Etc.



$uxv \quad uxxv$
 $uxv \quad \rightarrow uv$

Substitutable CF – summary

- repetition \leftrightarrow optionality
- one repetition \rightarrow indefinite repetition
- need $uX^nV, uX^{n+1}V$

Human language

- $HL \subseteq$ Mildly Context Sensitive Languages (Joshi, 1985)
- Not all optional elements in HL are repeatable, e.g. most Adverbs in English
- But HL is not describable by strict substitution classes (à la Zellig Harris) anyway
- Humans need additional info, e.g. meaning

- (4) a. The man in the yellow hat has a monkey
 b. He has a monkey
- (5) a. The man laughed
 b. The cheerful man laughed
- (6) a. The horse raced past the barn quickly.
 b. The horse raced past the barn fell.
 c. Sue fell.
 d. *Sue quickly

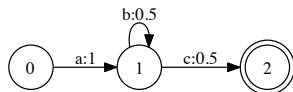
PDFa PAC learner

Clark and Thollard (2004)

- Learns Probabilistic Deterministic Finite State languages
- Requires huge sample sizes
- Similar to Angluin (1982)'s 0-reversible learner except states must share most of their suffixes (not just one) to be merged

PDFa PAC learner

Source grammar



Sample will be approximately:

50% *ac*

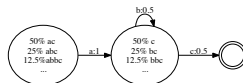
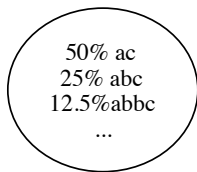
25% *abc*

12.5% *abbc*

6.25% *abbbc*

etc.

PDFA PAC learner



PDFA PAC learner: summary

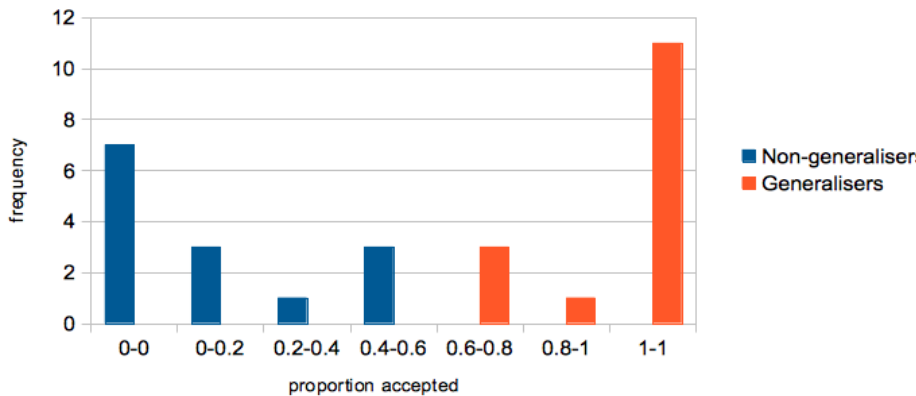
- Representative sample \rightarrow learns repetition
- Optionality $\not\leftrightarrow$ Repetition
- i.e. Learnable, but no shortcuts!

Artificial language learning

- **Experimental paradigm:** Expose people to a sample from a target language. See what they do with novel items from the language.
- **Experiment underway:**
 - Pilot method On-line survey (Amazon's Mechanical Turk)
 - Exposure stimuli *ac abc abbc* (a,b,c word classes)
 - Testing stimuli *ac abc abbc abbbc abbbbc*
 - Pilot results participants divide into two groups: generalisers (accept *abbbc*, *abbbbc*) and non-generalisers (reject *abbbc*, *abbbbc*)

Artificial language learning

Frequency of acceptance of generalised stimuli



Future research

- k, l -substitutable CFLs Yoshinaka (2008)
- CFLs with finite context and finite kernel properties Clark et al. (2008)
- k, l -substitutable Multiple Context Free Languages Yoshinaka (2009)
- What else? Accepting suggestions!
- Experiment

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n-gram learner

Just makes of list of the n-grams it has seen. String is grammatical if it contains no bigrams not in the list and has boundary markers.

Example (Bigram learner)

$\Sigma = \{a, b\}$, boundary markers \times , \times

Sample: a, ab, abb

Bigrams: $\times a, ab, bb, b\times, a\times$

Hypothesis language: ab^*

n-gram learner

Theorem

Let $u, v \in \Sigma^*$, $x \in \Sigma$. Let sample of L contain $uv, uxv, uxxv, \dots, ux^{n-1}v, ux^n v$. Then an n -gram learner will conclude that $ux^*v \subseteq L$.

Example (Trigram learner)

ababbbabab

$\times ab$ *abbbabab*
aba *bbbabab*
a *bab* *bbabab*
ab *abb* *babab*
aba *bbb* *abab*

abab *bba* *bab*
ababb *bab* *ab*
ababbb *aba* *b*
ababbbba *bab*
ababbbbab *ab* \times

Trigrams: $\times ab, aba, bab, abb, bbb, bba, ab \times$

n-gram learner

Hypothesis: $abab^*abab \subseteq L$

aba $bbbb$ $abab$

trigrams: $\times ab, aba, bab, abb, bbb, bba, ab \times$

$\times ab$ $abbbb$ $abab$
 aba $bbbb$ $abab$
 a bab bbb $abab$
 ab abb bb $abab$
 aba bbb b $abab$
 $abab$ bbb $abab$

$ababb$ bba bab
 $ababbb$ bab ab
 $ababbbb$ aba b
 $ababbbba$ bab
 $ababbbb$ $ab \times$

n-gram learner

If the repeated element is repeated fewer than n times, repetition isn't generalised.

Trigram learner Sample: *ababbabab*

\times ab *abbabab*
 aba *bbabab*
 a *bab* *babab*
 ab *abb* *abab*
 aba *bbb* *abab*

aba *bba* *bab*
 abab *bab* *ab*
 ababb *aba* *b*
 ababba *bab*
 ababbab *ab* \times

Trigrams: \times ab, aba, bab, abb, ~~bbb~~, bba, ab \times

* aba **bbb** bb...abab

n-gram learner summary

- Given up to n x 's in a row in context (u, v) , generalises to $ux^*v \subseteq L$
- No generalisation from optionality to repetition or vice-versa